An Invitation to Continuous Logic

Eric Tao University of Pennsylvania Directed Reading Program, December 2023 • A branch of mathematical logic

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- And yet has found applications in combinatorics, algebra, analysis, and more!
- This is a talk in *applied* model theory—I'll talk about a connection between model theory and functional analysis!

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Theorem

Let ϕ be a statement about fields. The following are equivalent:

- φ is true in every algebraically closed field of characteristic zero (one added to itself is never zero).
- 2. ϕ is true in some algebraically closed field of characteristic p (one added to itself p times is zero) for arbitrarily large p.

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Essence of the Proof

"Having zero characteristic" is not expressible as a single first-order statement!

The Ultraproduct!

$\left\{\overline{\mathbb{F}_p}\right\} \longrightarrow \mathbb{C}$

Fundamental Theorem of the Ultraproduct (Łoś)

For any discrete first-order sentence ϕ , ϕ true in most of the $\left\{\overline{\mathbb{F}_p}\right\} \Leftrightarrow \phi$ true in \mathbb{C}

But Not So Good For Complete Metric Structures...

$\{\mathbb{R}\} \longrightarrow \mathbb{R}$

Failure of Discrete First-order Logic

 \mathbb{R} is not complete!

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- 1968: Ax proves the Ax–Grothendieck Theorem using model theoretic methods.
- 2008: The de facto standard "Model Theory For Metric Structures" (Ben Yaacov, Berenstein, Henson, and Usvyatsov) for continuous logic is published, uniting the two ultraproducts.

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And all the results carry over! (Compactness, Löwenheim–Skolem, omitting types, elimination of imaginaries, stability theory, etc.)

The Connes Embedding Problem

Big open conjecture for a long time that was recently solved: Can I always embed a certain type of "operator algebra" into a certain ultrapower?

Equivalent to a number of different problems from all different fields across math and computer science:

- C*-algebras
- Quantum information theory
- Quantum complexity theory

And continuous logic forms the bridge!